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# A Fractal which violates the Axiom of Determinacy

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#### Abstract

By use of the axiom of choice I construct a symmetrical and self-similar subset  $A \subseteq [0,1] \subseteq \mathbb{R}$ . Then by an elementary strategy stealing argument it is shown that A is not determined. The (possible) existence of fractals like A clarifies the status of the controversial Axiom of Determinacy. <sup>1</sup>

In this note, I present an argument against the unrestricted axiom of determinacy (=AD).

Fix  $A \subseteq [0,1] \subseteq \mathbb{R}$ . We define an infinite game  $\mathcal{G}_A$  as follows. The initial position is I := [0,1]. Each position will be an interval  $[a,b] \subseteq [0,1]$ ,  $a,b \in \{\frac{p}{2^k}: p,k \in \mathbb{N}\}$ . In each position [a,b] there are always two legal moves. The player who has the turn can move "left" or can move "right". If the player moves "left" the new position is  $[a,\frac{a+b}{2}]$ . If the player moves "right" the new position is  $[\frac{a+b}{2},b]$ . Player  $\mathcal{A}$  makes the first move. In each actual game, successively the players construct a sequence  $I = [0,1] \supseteq I_1 \supseteq I_2 \supseteq \ldots$  of closed intervals. The interval  $I_{j+1}$  is either the left or the right closed interval of  $I_j$ . Each actual game produces a point  $p := \cap_{j \in \mathbb{N}} I_j \in [0,1]$ .

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According to the rules of  $\mathcal{G}_A$  player  $\mathcal{A}$  wins if  $p \in A$ . Otherwise (i.e. when  $p \notin A$ ) player  $\mathcal{B}$  wins.

A player has a winning strategy if there is a protocol (i.e. a map from the set of positions to { 'left', 'right'}) which guarantees victory. The set of positions can be divided into 3 classes. The positions which are won for player  $\mathcal{A}$ , the positions which are won for player  $\mathcal{B}$ , and the controversial class of the positions which are undetermined. According to the axiom of determinacy the last class is always empty. Each actual game has a winner. So if the players are clever enough it must be determined who will win the game. The intuition behind  $\mathbf{AD}$  is that if  $\mathcal{A}$  and  $\mathcal{B}$  have infinite powers the same player will win each game. If for instance  $\mathcal{A}$  wins the first game,  $\mathcal{A}$  ought also to win the second game. The argument is that if player  $\mathcal{B}$  wins this new game, sometimes in the first game  $\mathcal{B}$  could not have possibly played optimally  $^2$ .

It is well-known that  $\mathbf{AC}$  (the axiom of choice) and  $\mathbf{AD}$  are contradictory. The usual proof uses a diagonal argument combined with the fact that the number of strategies is  $2^{\aleph_0}$  [1], [2]. The status of  $\mathbf{AD}$  has been examined in great depth [4],[5]. There seems to be two approaches. One can accept  $\mathbf{AC}$  and ask which sets are determined. This leads to questions which are independent of the usual axiomatization of set theory [2],[4], [5]. The other and more radical approach is to discard the axiom of choice [3]. The main argument is that  $\mathbf{AD}$  is deductively strong and has many nice consequences, [2],[4]. Still there is no doubt that most of us prefer  $\mathbf{AC}$ .

Let  $\mathbb{Q}$  denote the rational numbers.

**Theorem (AC)** There exists a set  $A \subseteq [0,1]$  such that A and  $A \cap [0,\frac{1}{2}]$  are isomorphic under the map  $x \to \frac{x}{2}$ , and  $A \setminus \mathbb{Q}$  and  $A^c \setminus \mathbb{Q}$  are isomorphic under the map  $x \to 1-x$ .

**Proof:** Consider the collection **J** of pairs (A, B) where  $A, B \subseteq [0, 1] \setminus \mathbb{Q}$ , where  $A \cap B = \emptyset$  and where  $A = A \cdot \mathbb{Q} \cap [0, 1] \setminus \mathbb{Q}$  and  $B = B \cdot \mathbb{Q} \cap [0, 1] \setminus \mathbb{Q}$ . The set **J** of such pairs are ordered inductively under inclusion. According to Zorn's lemma (which is equivalent to **AC**) there must be a pair  $(A, B) \in \mathbf{J}$  which is maximal with respect to inclusion. We claim that  $A \cup B = [0, 1] \setminus \mathbb{Q}$ . Otherwise there would be  $x \in ([0, 1] \setminus \mathbb{Q}) \setminus (A \cup B)$ . Notice that,  $(1-x) \cdot \mathbb{Q} \cap A = \emptyset$ ,  $x \cdot \mathbb{Q} \cap (1-x) \cdot \mathbb{Q} = \emptyset$  and  $x \cdot \mathbb{Q} \cap B = \emptyset$ . Thus  $(A \cup x \cdot \mathbb{Q}, B \cup (1-x) \cdot \mathbb{Q})$ 

 $<sup>^2</sup>$ Unless of course  $\mathcal{A}$  first deviate from the line of play in the first play. But to deviate and lose does not seem wise.

is well-defined and belongs to **J**. This violates the maximality of (A, B).  $\square$  **Theorem** All the positions  $(I', A \cap I')$  are isomorphic, when the points in  $\mathbb{Q}$  are ignored. No move can make any difference to the outcome because all positions are isomorphic. Each game produces a winner.

**Proof:** We ignore the points in  $\mathbb{Q}$ . Notice that the two positions which can be reached from I both are isomorphic to I (when the role of A is changed with that of  $A^c$ ). This is because  $f_1: x \to 1-2x$  maps  $A^c \cap [0, \frac{1}{2}]$  isomorphic onto A, and because  $f_2: x \to 2-2x$  maps  $A^c \cap [\frac{1}{2}, 1]$  isomorphic onto A.  $\square$  Corollary (AC)  $\mathcal{G}_A$  is not determined.

**Proof:** If player  $\mathcal{A}$  has a winning strategy, player  $\mathcal{B}$  can steal it. It is not difficult to show that the points in  $\mathbb{Q}$  do not affect this argument.

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