

Polymorphic Embedding of DSLs

based on Hofer, Ostermann, Rendel, Moors, *Polymorphic Embedding of DSLs*, ACM Conference on Generative Programming and Component Engineering, 2008

How does this fit into dOvs?

- EDSLs and compilers have a lot in common:
 - translation of higher-level entities into lower-level entities
 - checking well-formedness
 - optimizations
- Main difference: EDSLs are compositional by construction (which is both good and bad)
- This approach shows how to stick with the basic EDSL approach while making it almost as flexible as a compiler
- It can also be seen as a tutorial on how to use compiler techniques in libraries.
- You'll learn some Scala on the way, too.

The Traditional Approach

- P. Hudak, Modular Domain Specific Languages and Tools
- DSL as **library**, not as a separate language
- DSL as an **algebra**, not via building ASTs
- Example: A **Regions** language

```
type Region = Vector ⇒ boolean
```

```
def univ : Region = p ⇒ true
```

```
def circle : Region
```

```
  = p ⇒ p._1 * p._1 + p._2 * p._2 < 1
```

```
def union(x : Region, y : Region) : Region
```

```
  = p ⇒ x(p) || y(p)
```

```
...
```

Pros and Cons

- Pros
 - **Reuse** of language infrastructure (incl. **type checking**)
 - Interpretation is **compositional** (defined by an algebra)
 - Allows combining several DSLs
- Cons
 - The interpretation is integral part of the language
 - **Alternative interpretations** cannot be supplied
 - Interpretations are not **components**
 - In particular: **Optimizations** cannot be applied to them

Contributions

- Pure Embedding with **multiple** interpretations
 - Analyses and optimizations as “yet another” interpretation
- Interpretations and languages as **components**
- Scala as implementation language in OO context
 - Show-case for advanced OO features (partly experimental)

Explicit language interface

```
trait Regions {  
  // Ordinary type synonyms  
  type Vector = (double, double)  
  
  // Abstract domain types  
  type Region  
  
  // Abstract domain operations  
  def univ : Region  
  def circle : Region  
  def union(x : Region, y : Region) : Region  
  def scale(v : Vector, r : Region) : Region  
  ...  
}
```

Explicit language interface II

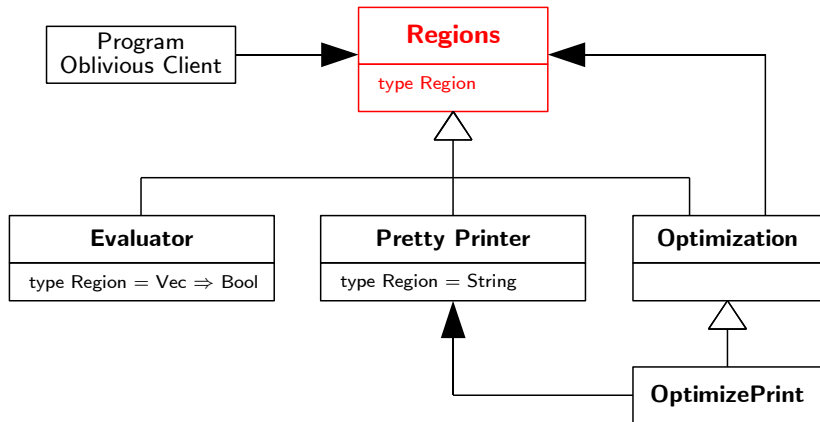
- **Abstract type members** represent domain types

```
type Region
```

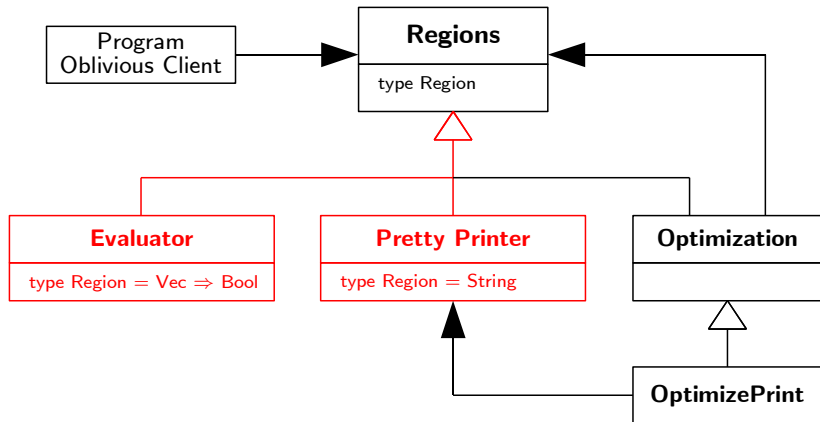
- Compositional by construction:
Interface is the **signature of the algebra**

```
def union(x : Region, y : Region) : Region
```

Architecture Overview



Architecture Overview



An Evaluator

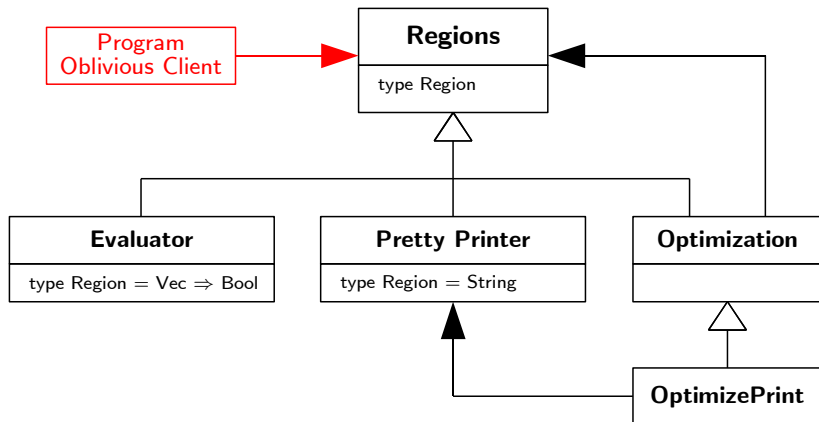
```
trait Evaluation extends Regions {  
  type Region = Vector => boolean  
  
  def univ : Region = p => true  
  def circle : Region  
    = p => p._1 * p._1 + p._2 * p._2 < 1  
  def union(x : Region, y : Region) : Region  
    = p => x(p) || y(p)  
  ...  
}  
object Eval extends Evaluation
```

- Same definitions as in the traditional approach

A Pretty Printer

```
trait Printing extends Regions {  
  type Region = String  
  
  def univ : Region = "univ"  
  def circle : Region = "circle"  
  def union(x : Region, y : Region) : Region  
    = "union(" + x + ", " + y + ")"  
  ...  
}  
object Print extends Printing
```

Architecture Overview

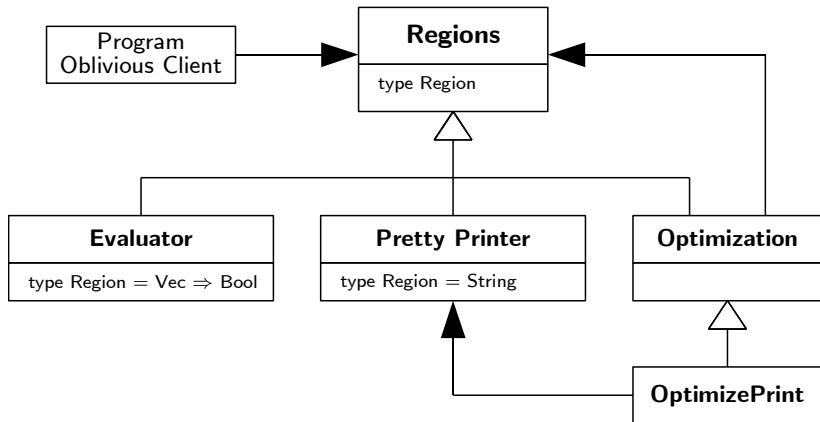


Programs as Oblivious Clients

```
// A simple program
def program( semantics : Regions )
    : semantics.Region = {
    import semantics._
    val ellipse24 = scale((2, 4), circle)
    union(univ, ellipse24)
}
```

- A DSL program has **path-dependent type**: `semantics.Region`
- `println(program(Eval)((1, 2)))` prints **true**
- `println(program(Print))` prints **union(univ, scale((2, 4), circle))**

Architecture Overview



A DSL with Polymorphism

- Example: Functions language (inspired by Carette et al.)
- User-defined bindings

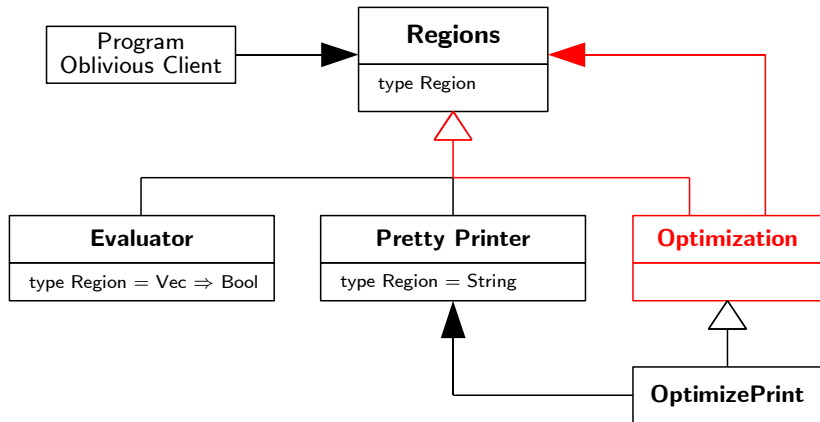
```
trait Functions {  
  // Abstract domain types  
  type Rep[X]  
  // Abstract domain operations  
  def fun[S, T](f : Rep[S] => Rep[T])  
    : Rep[S => T]  
  def app[S, T](f : Rep[S => T], v : Rep[S])  
    : Rep[T]  
}
```

- Using **higher-kinded abstract type member** Rep
- Using higher-order abstract syntax (**HOAS**)

Two Example Interpretations

```
trait FunEval extends Functions {  
  type Rep[T] = T  
  def fun [S,T](f : S ⇒ T) = f  
  def app [S,T](f : S ⇒ T, v : S) : T = f(v)  
}  
  
trait FunPrinting extends Functions {  
  type Rep[X] = String  
  def fun [S,T](f : String ⇒ String) : String  
    =  
    {  
      val v = variables.next  
      "fun(" + v + " ⇒ " + f(v) + ")"  
    }  
  ...  
}
```

Architecture Overview

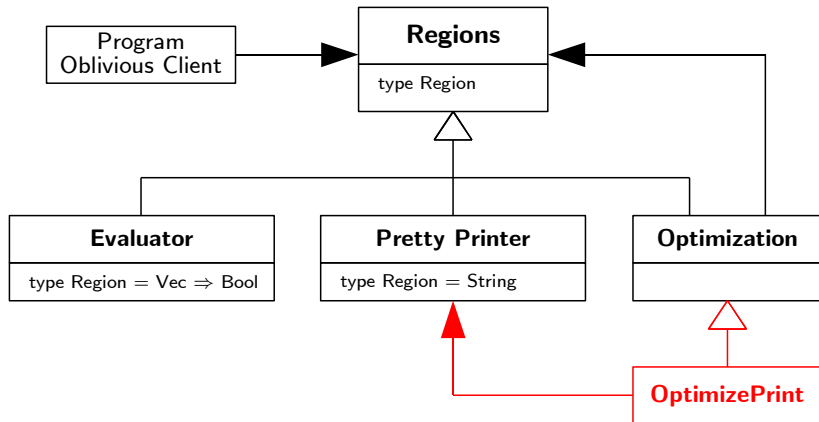


Interpretations as Components

- Example: Optimization

```
trait Optimization extends Regions {  
  val semantics : Regions  
  
  type Region = (semantics.Region, boolean)  
  def univ : Region = (semantics.univ, true)  
  def circle : Region =  
    (semantics.circle, false)  
  def union(x : Region, y : Region) : Region  
    =  
    if (x._2 || y._2) (semantics.univ, true)  
    else (semantics.union(x._1, y._1), false)  
  ...  
}
```

Architecture Overview



Reuse of Interpretations

- Interpretations can be regarded as **reusable components**
 - Odersky / Zenger: Scalable Component Abstractions
- Example: An **optimizing** interpretation can work on several interpretations

```
object OptimizePrint extends Optimization {  
  val semantics = Print  
}
```

- `println(program(OptimizePrint))` prints
(univ, true)
- `while println(program(Print))` prints
union(univ, scale((2, 4), circle))

Hierarchical Composition

- Example: A Vectors **sublanguage**

```
trait Vectors {  
  type Vector  
  ...  
}
```

```
trait Regions {  
  val vec : Vectors  
  import vec._  
  ...  
}
```

Interplay with Interpretation Components

- Example: Optimization for refactored Regions language
- Needs **singleton types**

```
trait Optimization extends Regions {  
  val semantics : Regions  
  val vec : semantics.vec.type = semantics.vec  
  import vec._  
  ...  
}
```

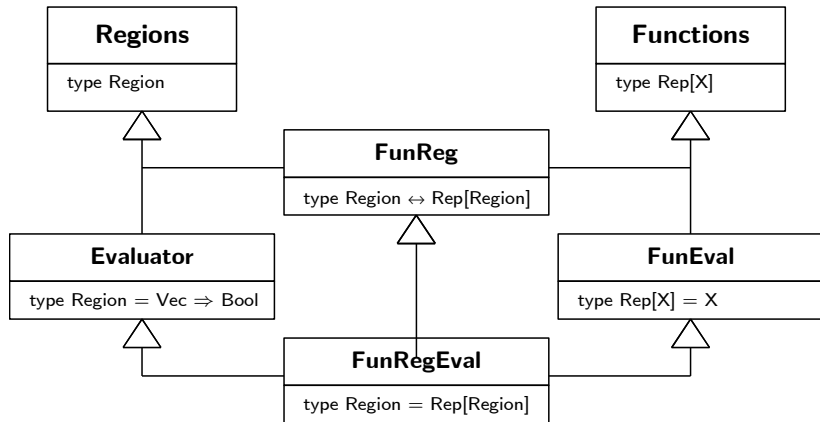
Peer Composition

- Example: Combine Regions with Functions language
- Problem: Representations have to be translated

```
trait FunReg extends
    Regions with Functions {
  implicit def fromRegion(r : Region)
    : Rep[Region]
  implicit def toRegion(r : Rep[Region])
    : Region
}
```

- Using Scala's *implicit* conversions for less verbosity

Peer Composition: Overview



Peer Composition of Interpretations

- Integration for both evaluation and printing semantics
- Example: Evaluation

```
object FunRegEval extends FunReg
  with Evaluation with FunEval {
  implicit def fromRegion(r : Region)
    : Rep[Region] = r
  implicit def toRegion(r : Rep[Region])
    : Region = r
}
```

Summary

- Reuse of the language infrastructure in **pure embedding** style
- Interpretation components
 - In particular: Application of **optimizations** on them
- Language components

- Outlook
 - Compositionality can be limiting
 - Regard Scala arithmetics, etc. as language interfaces
 - Alternative approaches
 - Type classes (Haskell)
 - Virtual classes (gbeta)